### BITS Pilani, Dubai Campus Dubai International Academic City, Dubai 1 Semester 2010-11

Course

**CS C351 Theory of Computation** 

Component

Comprehensive Examination

Year

III Year CS

Date

30-12-2010

Duration

3 hrs

Weightage

40% (80 Marks)

No. of Pages

4 Pages

No. of Section

A&B

Note:- Answer the Section A answers in Blue Answer book and the Section B in Green Answer Book.

## Section A

1. What does the following Turing machine transition function meant for? Justify (5 Marks) and prove your answer using a sample input string.

Current symbol	Write	Move	New state
*	*	Right	GET1
0	0	None	HALT
1	X	Right	DROP1
у	1	Right	CHANGEY
0	у	Left	FINDX
1	1	Right	DROP1
y y	у	Right	DROP1
1	1	Left	FINDX
у	у	Left	FINDX
х	х	Right	GET1
0	0	Left	CHANGEX
У	1	Ríght	CHANGEY
1	1	Left	CHANGEX
X	1	Left	CHANGEX
*	*	None	HALT
	*  0  1  y  0  1  y  1  y  1  y  1  x	* *  0 0  1	*       *       Right         0       0       None         1       X       Right         y       1       Right         0       y       Left         1       1       Right         y       y       Right         1       1       Left         y       1       Right         0       0       Left         y       1       Right         1       1       Left         x       1       Left         Left       Left       Left

2. Design a finite automaton for the language given below:

L = {w | w contains equal number 0's and 1's}.

Your design should include the state called for **reject state** to terminate the computation process if the input string does not satisfy the above language.

(7.5 Marks)

3. Convert the regular expression into finite automaton.

(7.5Marks)

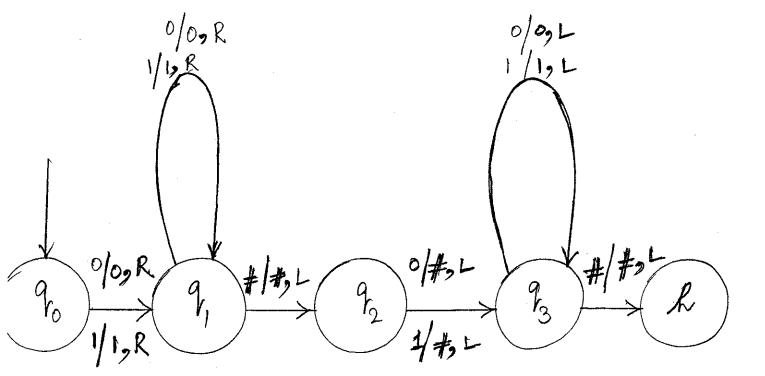
(0 + 11 + 01)\*0\*(01)\*

- 4. Design a finite automata and a Turing Machine that reads binary strings and performs the following actions. If the input represents an odd number, subtracts one to the number. If the input represents an even number, add one to the number. For example, for input "101" the output should be "100" and for input "1010" the output should be "1011". (10 Marks)
- 5. What do you mean by P and NP Completeness?

(3 Marks)

6. What does the following transition diagram represent? Justify your answer with proper configuration history. (7 Marks)

Where P is the halt state.



#### Section B

1. Convert the following CFG to PDA.

S -> aAA

A -> aS | bS | a

Give the definition and transition functions of PDA. Check with a string. (5 Marks)

2. Let  $G = (V, \sum, R, Expr)$  be a CFG with variables

V = { Expr, Factor, Term}, and terminals.

 $\Sigma$  = { CONSTANT, IDENTIFIER, PLUS, TIMES, LPAREN, RPAREN} and rules:

Expr -> Term | Expr PLUS Term

Term -> Factor | Term Times Factor

Factor -> IDENTIFIER | CONSTANT | LPAREN Expr RPAREN

Here are regular expression for the terminals:

CONSTANT = (0 U 1 U...U 9) +

IDENTIFIER = (a U b U...U z U A U B U... U Z) \*

PLUS = +

TIMES = \*

LPAREN = (

RPAREN = )

White space between terminals is ignored.

For each string below, show that it is generated by G by showing a derivation or explain why it is not generated by G. (3\*5=15Marks)

- a. 2\*a+b
- b. (aardvark+2)\*antelope
- c. 2x+3\*(y+z)
- 3. Specify CFG for the language.

(10 Marks)

 $L = \{x \in \{a,b,c\}^* : x \text{ contains an equal number of two different characters}\}$ 

4. Let  $G = (V, \sum, R, Expr)$  be the following grammar:

(5 Marks)

STMT → ASSIGN | If Then | If Then Else

If Then → if condition then STMT

If Then Else→ if condition then STMT else stmt

ASSIGN → a:= 1

 $\sum$  = {if, condition then, else, a:=1}

V = {STMT, If Then, If Then Else, ASSIGN}

Check whether the grammar is ambiguous or not for the sentence "if condition then if condition then a:= 1 else a:=1.

5. PDA P = ( $\{q0,q1,q2,q3,f\}, \{a,b\}, \{Z0,A,B\}, \delta, q0, Z0, \{f\}$ ) has the following rules defining  $\delta$ . (5 Marks)

$$\delta(q0, a, Z0) = (q1, AAZ0)$$
 $\delta(q0, b, Z0) = (q2, BZ0)$ 
 $\delta(q0, \epsilon, Z0) = (f, \epsilon)$ 
 $\delta(q1, a, A) = (q1, AAA)$ 
 $\delta(q1, b, A) = (q1, \epsilon)$ 
 $\delta(q1, \epsilon, Z0) = (q0, Z0)$ 
 $\delta(q2, a, B) = (q3, \epsilon)$ 
 $\delta(q2, b, B) = (q2, BB)$ 
 $\delta(q2, \epsilon, Z0) = (q0, Z0)$ 
 $\delta(q3, \epsilon, B) = (q2, \epsilon)$ 
 $\delta(q3, \epsilon, Z0) = (q1, AZ0)$ 

Give an execution trace (sequence of ID's ) showing that string bab is in L(P).

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#### BITS, PILANI, DUBAI DUBAI INTERNATIONAL ACADEMIC CITY, DUBAI I-SEMESTER 2010-11

Course

CS C351 Theory of Computation

Component

Test - 2 (Open Book)

Duration

50 mins

Date

12-12-2010

Weightage (%)

20 % (40 Marks)

Year

III Year Computer Science

No. of Pages

02 Pages

Note: Make necessary assumption wherever it is required. Start answering the answers in a fresh page.

Only Text Book, Class notes and Photocopy of the Turing Machine slides are allowed.

1. Construct a PDA for the CFG given below:

(10 Marks)

 $| \rightarrow a | b | Ia | Ib | IO | I1$ 

 $E \rightarrow ||E^*E||E+E||(E)$ 

2. For the given CFG grammar trace the Shift-Reduce parsing of the input string: x - 2 \* y \$.

(10 Marks)

S = E\$

 $E = E+T \mid E-T \mid T$ 

T = T\*F | T/F | F

 $F = num \mid id$ 

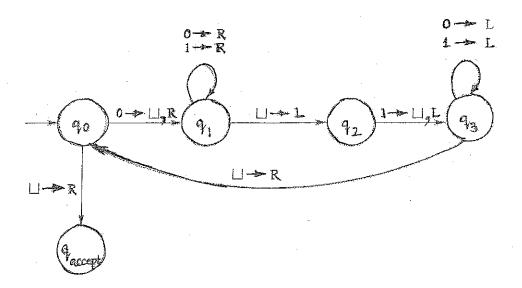
3. Design a Turing Machine to check well balanced structure of Parenthesis, using the set of transition function defined below for the input string \$ ( ( ( ) ) ( ) ) \$. Give the Transition diagram and the table. \$ represent the blank space. (10 Marks)

$$(J0, () = (J1,X,R); (J0,Y) = (J0,Y,R); (J0,$) = (J3,$L);$$

$$(J1, () = (J1, (, R); (J1, )) = (J2, Y, L); (J1, Y) = (J1, Y, R);$$

$$(J2, ()) = (J2, (, L); (J2, X) = (J0, X, R); (J2, Y) = (J2, Y, L);$$

$$(J3, X) = (J3, X, L); (J3, Y) = (J3, Y, L); (J3, $) = (Accept)$$



## BITS, Pilani-Dubai Dubai International Academic City, Dubai I-Semester 2010-11

Course Number

CS C351

Course Name

: Theory of Computation

Year

: III Year CS

Component

: Test - 1 (Closed Book)

Weightage / marks

25 % / 50 Marks

Date

: 31-10-2010

## Note: answer all the questions.

1. Find a derivation tree of a \* b + a \* b given that a \* b + a \* b is in L(G), where G is given by  $S \rightarrow S + S \mid S * S$ ,  $S \rightarrow a \mid b$ .

2. In both cases below, a transition table is given for a PDA with initial state q0 and accepting state q2. Describe in each case the language that is accepted.

Move Number	State	Input	Stack Symbol	Move(s)
1	q0	a	Z0	(q1, aZ0)
2	<b>q</b> 0	ь	<b>Z</b> 0	(q1, bZ0)
3	ql	a	a	(q1, a), (q2,a)
4	q1	b	a	(q1,a)
5	q1	a	b	(q1,b)
6	ql	b	ь	(q1,b), (q2,b)

Move Number	State	Input	Stack Symbol	Move(s)
1	q0	a	Z0	(q0, XZ0)
2	q0	ь	<b>Z</b> 0	(q0, XZ0)
3	q0	a	X	(q0, XX)
4	q0	b	X	(q0, XX)
5	q0	C	X	(q1, X)
6	q0	С	Z0	(q1, Z0)
7	q1	a	X	(q1, null)
8	q1	b	X	(q1, null)
9	q1 :	nul1	Z0 ·	(q2, Z0)

3. Check whether the given Context-Free Grammar is ambiguous or not for the input string  $a^2b^2a^2$ .

(5Mooks)

Production rules are:  $S \rightarrow aAS \mid a$ ,

$$A \rightarrow SbA \mid SS \mid ba$$
.

- 4. Write a deterministic finite automaton to recognize each of the following languages.
  - a.  $L1 = \{w \in \{a,b\} * | \text{ each } a \text{ in } w \text{ is immediately preceded and immediately followed by a } b\}$ (Use 5 states)
  - b. L2 = {  $w \in \{a,b\} * | w \text{ has both } ab \text{ and } ba \text{ as substrings}\}.(Use 6 \text{ states})$

$$(2 * 5 = 10 \text{ marks})$$

5. Design a nondeterministic finite automaton for the following language. (5 marks)

$$L = \{ x \in \{0,1\} * | x \text{ has } 101 \text{ as a substring.} (Use 4 \text{ states}) \}$$

- 6. Write a Context Free Grammar for the following languages.
- (2 \* 5 = 10 marks)
- a. L accepts strings on the alphabet {a, b, c}, where number of c's equals the sum of the number of 'a's and 'b's.
- b.  $L = \{ a^n b^m | 2n \le m \le 3n \}$

BITS. Pilani - Dubai

Course:	Theory,	Dubai Internat I – S Computation	ional Academic Cit Semester 2010-11 Quiz – II		II year	CS
ld. No:			Name:			_
Section	n:		_ Duration: 20	) mins		

1. Construct a CFG grammar accepting  $\{a^nb^ma^n \mid m, n >= 1\}$  by null store. The transition functions are  $\delta(q_0, a, z_0) = \{ (q_0, aZ_0) \}$ 

Weightage: 7 % (Marks: 14 Marks)

 $\delta (q_0, a, a) = \{(q_0, a_0)\}$ 

 $\delta (q_0, b, a) = \{(q_1, a)\}$ 

 $\delta (q_1, b, a) = \{(q_1, a)\}$ 

 $\delta$  (q<sub>1</sub>, a, a) = {(q1, null)}

 $\delta$  (q<sub>1</sub>, null, Z<sub>0</sub>) = {(q1, null)}

2. Construct a PDA for the CFG given below for the language L of all balanced strings involving two types of brackets, "{ }" and "[ ]". The language L is generated by the CFG with production:

 $S \rightarrow SS \mid [S] \mid \{S\} \mid null$ 

Answers:

Date: 22-11-2010

# BITS, PILANI-DUBAI DUBAI INTERNATIONAL ACADEMIC CITY, DUBAI I- SEMESTER 2010-11



Name:		Date: 06 \$	Sept 2010
ld. No.:		Section: _	/
Course:	CS C351 Theory of Computation	III Year CS Qu	IZ-I
Marks: 16	Marks ( 8 %)	t ·	
1. [	Draw Deterministic Finite Automata to accep	ot the following sets of strings o	ver the
а	lphabet {0,1}:	4	Marks
	(i) All strings ending in "1101" using 5	states.	

(ii) All strings that start with 0 and has odd length or start with 1 and has even length, using 3 states.

2.	All strings th	nat contains	an even	number	of Os	or exactly two	1s.	(6 states)
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4 Marks

3.	Generate the given (a1*b2+a1)*(a1+b2) using the CFG Grammar p	roduc	ction rules
	and also give the CFG definition for this problem.	4	Marsky

Production rules are:

E	$\rightarrow$	E-	+7	• •		٠.	٠.	٠.	٠.	 	 	• •	 • •	 	 ٠.	 	•	 	٠.	 	 	R	lul	е	1	
Ε	$\rightarrow$	Т			٠	٠.	٠.		•••	 	 			 	 	 				 	 		R۱	ule	∋ 2	2
_	,	<b>-</b> :	* r	_																			_	_		

T → F
F → (E) Rule 5
F → a1Rule 6
F → b2Rule 7